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## Path Cover Number of Rook's graph

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### ABSTRACT

The path cover number of graph  $G$ , is a set of induced paths of  $G$  that cover all vertices of graph  $G$ . The path cover number  $P(G)$  of graph  $G$  is the smallest positive integer  $k$  such that there are  $k$  vertex-disjoint path occurring as induced subgraphs of  $G$  that cover all the vertices of  $G$ . There are some relationships between  $P(G)$  and some other graph parameters such as minimum rank of graphs and zero forcing number, etc.

Here we investigate the path cover number of Rook's graph and vertex-sum of two Rook's graphs.

**KEYWORDS:** path covering, path cover number, vertex-sum

### 1 INTRODUCTION

Let  $G = (V_G, E_G)$  be a graph with vertex set  $V_G$  and edge set  $E_G$ . The order of  $G$  is denoted by  $|G|$  which is equal to cardinality of  $V_G$ . Throughout this paper, all graphs are simple (no loops, no multiple edges), undirected, and have finite nonempty vertex sets. The cartesian product  $G \square H$  of two graphs  $G$  and  $H$  has vertex set  $V_G \times V_H$  and two vertices  $(g_1, h_1)$  and  $(g_2, h_2)$  are adjacent if and only if either  $g_1 = g_2$  and  $h_1 h_2 \in E_H$ , or  $h_1 = h_2$  and  $g_1 g_2 \in E_G$ .

Path covering of a graph  $G$  is a family of induced disjoint paths in  $G$  that cover all the vertices of the graph. The minimum number of such paths is the path cover number of  $G$  and is denoted by  $P(G)$ .  $P(G)$  was first used in the study of minimum rank and maximum eigenvalue multiplicity in [1]. It is also related to some other graph parameters like zero forcing number  $Z(G)$  which is the minimum size of a zero forcing set.

Let  $G$  be a graph with each vertex given either colour black or white. The colour change rule is: If  $u$  is a black vertex and  $v$  is the only white neighbor of  $u$ , then change colour of  $v$  to black and in this case we say  $u$  forces  $v$  and write  $u \rightarrow v$ . In [3], the authors showed that for any graph the zero forcing number is an upper bound for the path cover number (i.e., for any graph  $G$ ,  $P(G) \leq Z(G)$ ).

## 2 ROOK'S GRAPH

A graph can be formed from an  $m \times n$  chessboard by taking the squares as the vertices and two vertices are adjacent if a chess piece situated on one square covered the other. For example Queen's graph, Rook's graph, etc. The Rook's graph  $R_{mn}$  is the graph that describes all possible movements of a rook as a chess piece on an empty  $m \times n$  chessboard. Rooks move either horizontally or vertically. Thus, the Rook's graph can represent the graph  $K_n \square K_m$  which is the cartesian product of two complete graphs. A square  $(i,j)$  indicates the vertex located on the  $i$ -th copy of  $K_n$  and the  $j$ -th copy of  $K_m$  at the same time. All the squares of row  $i$  and column  $j$  are the neighbors of the square  $(i,j)$ .

**Definition 2.1.** [2] A  $k$ -colouring of a graph  $G$  is a labeling  $f : V(G) \rightarrow S$ , where  $|S|=k$  (often we use  $S=[k]$ ). The labels are colours; the vertices of one colour form a colour class. A  $k$ -colouring is proper if adjacent vertices have different labels. A graph is  $k$ -colourable if it has a proper  $k$ -colouring. The chromatic number  $\chi(G)$  is the least  $k$  such that  $G$  is  $k$ -colourable.

A Latin rectangle is an  $m \times n$  matrix ( $m \leq n$ ) that has the numbers  $1, 2, 3, \dots, n$  as its entries with no number occurring more than once in any line (row or column).

Comparing previous definitions gives us the following proposition.

**Proposition 2.2.** Every  $m \times n$  latin rectangle represents a proper colouring of  $R_{mn}$  which is the graph  $K_n \square K_m$ .

*Proof.* It is obvious that every colour class is a (broken) diagonal. So every homochromatic vertices are non-adjacent and the colouring is proper.

Now, we try to introduce a new colouring for Rook's graph such that every colour class is an induced path of the graph.

Consider the graph  $R_{mn}$  when  $m < n$  and  $n$  is even. First place two  $k$  for  $1 \leq k \leq \frac{n}{2}$  in the squares  $(1,2k-1)$  and  $(1,2k)$ . Then place the other "k"s in the successive squares in a stairway order along two parallel diagonal which slopes downward and to the right, with the following modifications:

1. When the bottom row for a specific  $k$  is reached, stop the diagonal process for this  $k$  in this diagonal.

2. When the right-hand column is reached, the next  $k$  is put in the left-hand column as if immediately succeeded the right-hand column and continue the diagonal process to reach the bottom row.

Every diagonal has  $m$  squares hence every integer appears  $2m$  times in this graph. Each colour class forms an induced path of graph  $R_{mn}$ . Every square has exactly one integer so these paths are vertex-disjoint. Therefore we can present following lemma:

**Lemma 2.3.** For two integers  $m, n$  when  $m < n$  and  $n$  is even, we have:

$$P(R_{mn}) \geq \frac{n}{2}.$$

We illustrate this path covering in figure 1 for graph  $K_{10} \square K_5$  ( which is the graph  $R_{5,10}$ ).

1	1	2	2	3	3	4	4	5	5
5	1	1	2	2	3	3	4	4	5
5	5	1	1	2	2	3	3	4	4
4	5	5	1	1	2	2	3	3	4
4	4	5	5	1	1	2	2	3	3

Figure 1: The corresponding  $5 \times 10$  chessboard of graph  $K_{10} \square K_5$  with the squares of colour  $k$  representing the  $k$ -th induced path.

We can prove that the equality always holds.

**Theorem 2.4.** For two integers  $m, n$  when  $m < n$  and  $n$  is even, we have:

$$P(R_{mn}) = \frac{n}{2}.$$

*Proof.* The upper bound comes from lemma 2.3. It is enough to prove that  $P(R_{mn}) > \frac{n}{2}$ . Assume, for the sake of contradiction, that  $P(R_{mn}) \geq \frac{n}{2} - 1$ . There are  $mn$  squares in  $m$  rows. The pigeonhole principle then asserts that there is a path in this covering that has at least three squares in one row. These three squares induced a cycle which is a contradiction. So every induced path in this graph has at most  $2m$  vertices. Then  $P(R_{mn}) \geq \frac{n}{2}$  and the equality holds.

The authors guess that for other integer numbers  $m$  and  $n$ , the following results hold:

**Theorem 2.5.** For every integer  $n \leq 3$ ,

$$P(R_{nn}) = \begin{cases} n & n \text{ is odd} \\ n - 1 & n \text{ is even} \end{cases}$$

**Theorem 2.6.** For two integers  $m, n$  when  $m < n$  and  $n$  is odd, we have:

$$P(R_{mn}) = \left\lfloor \frac{m+n}{2} \right\rfloor.$$

### 3 VERTEX-SUM OF TWO ROOK'S GRAPHS

Let  $G$  and  $H$  be two vertex-disjoint graphs and assume that  $v_1$  and  $v_2$  are two vertices of  $G$  and  $H$  respectively, then the vertex-sum of  $G$  and  $H$  over  $v_1$  and  $v_2$  is the graph formed by identifying  $v_1$  and  $v_2$  to a unique vertex  $v$ , which is denoted by  $G +_v H$ .

**Theorem 3.1.** [4] Let  $G$  and  $H$  be two graphs each with a vertex labeled  $v$ . Then the following hold:

1. If there is a minimal path covering of  $G$  in which  $v$  is a path of length 1 and  $v$  is not an end-point in any minimal path covering of  $H$ , then

$$P(G +_v H) = P(G) + P(H) - 1.$$

∇. If  $v$  is an end-point of some minimal path covering of  $G$  but no minimal path covering contains  $v$  as a path of length 1 and  $v$  is not an end-point in any minimal path covering of  $H$ , then

$$P(G +_v H) = P(G) + P(H).$$

∕. If there are minimal path coverings of  $H$  and  $G$  in which  $v$  is an end-point of a path in both, then

$$P(G +_v H) = P(G) + P(H) - 1.$$

ξ. If  $v$  is neither an end-point of any minimal path covering of  $H$  nor the end-point in any minimal path covering of  $G$ , then

$$P(G +_v H) = P(G) + P(H) + 1.$$

In vertex-sum of two Rook's graphs, the locations of vertices  $v_1$  and  $v_2$  is not important because we can relabeling the vertices such that  $v_1$  is the square  $(m,n)$  in  $G$  and  $v_2$  is the square  $(1,1)$  in  $H$ . Then  $v_1$  and  $v_2$  are end-points in the minimal path covering which we introduced before lemma 2.3. Hence we have the following corollary.

**Corollary 3.2.** For every integers  $m, n, r$  and  $s$ ,

$$P(R_{mn} \square R_{rs}) = P(R_{mn}) + P(R_{rs}) - 1.$$

Especially when  $m < n$ ,  $r < s$  and  $n, s$  are even, we have:

$$P(R_{mn} \square R_{rs}) = \left\lfloor \frac{n+s}{2} \right\rfloor.$$

**Example 3.3.** A minimal path covering of graph  $(K_8 \square K_5) +_v (K_6 \square K_4)$  is:

1	1	2	2	3	3	4	4						
4	1	1	2	2	3	3	4						
4	4	1	1	2	2	3	3						
3	4	4	1	1	2	2	3						
3	3	4	4	1	1	2	2	2	5	5	6	6	
								6	2	2	5	5	6
								6	6	2	2	5	5
								5	6	6	2	2	5

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